

EECS150 - Digital Design
Lecture 17 – Memory 2

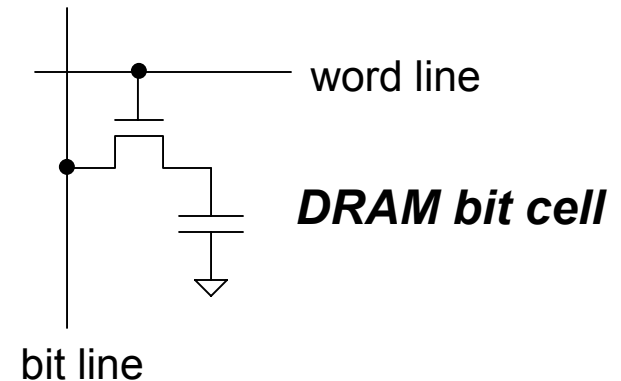
October 22, 2002

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SDRAM Recap

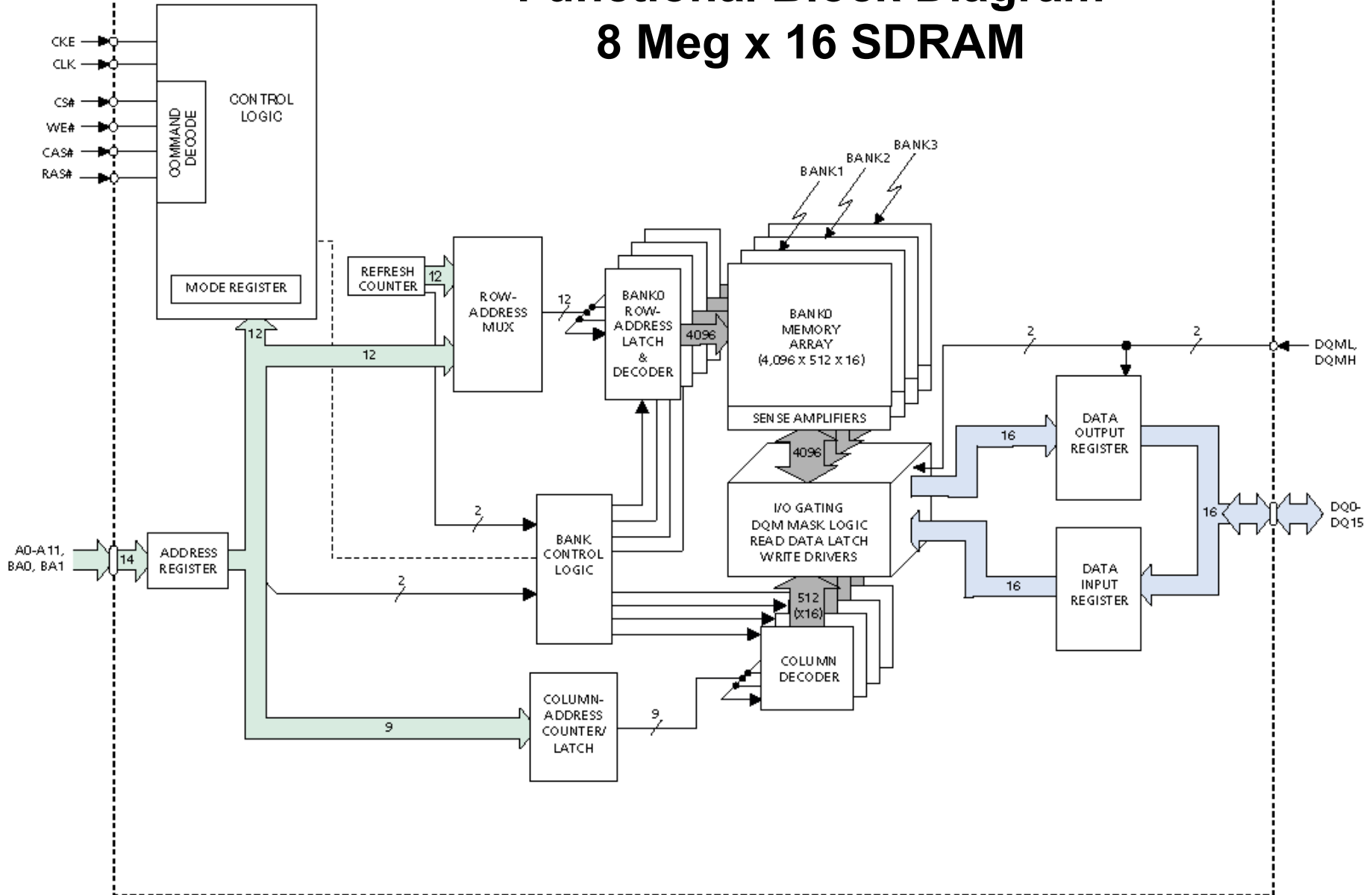
- General Characteristics

- Optimized for high density and therefore low cost/bit
- Special fabrication process – usually on a separate chip from processor
- Needs periodic refresh (in most applications)
- Relatively slow because:
 - High capacity leads to large cell arrays with high word- and bit-line capacitance
 - Complex read/write cycle. Read needs “precharge” and write-back



- Multiple clock cycles per read or write access
- Multiple reads and writes are often grouped together to amortize overhead. Referred to as “bursting”.

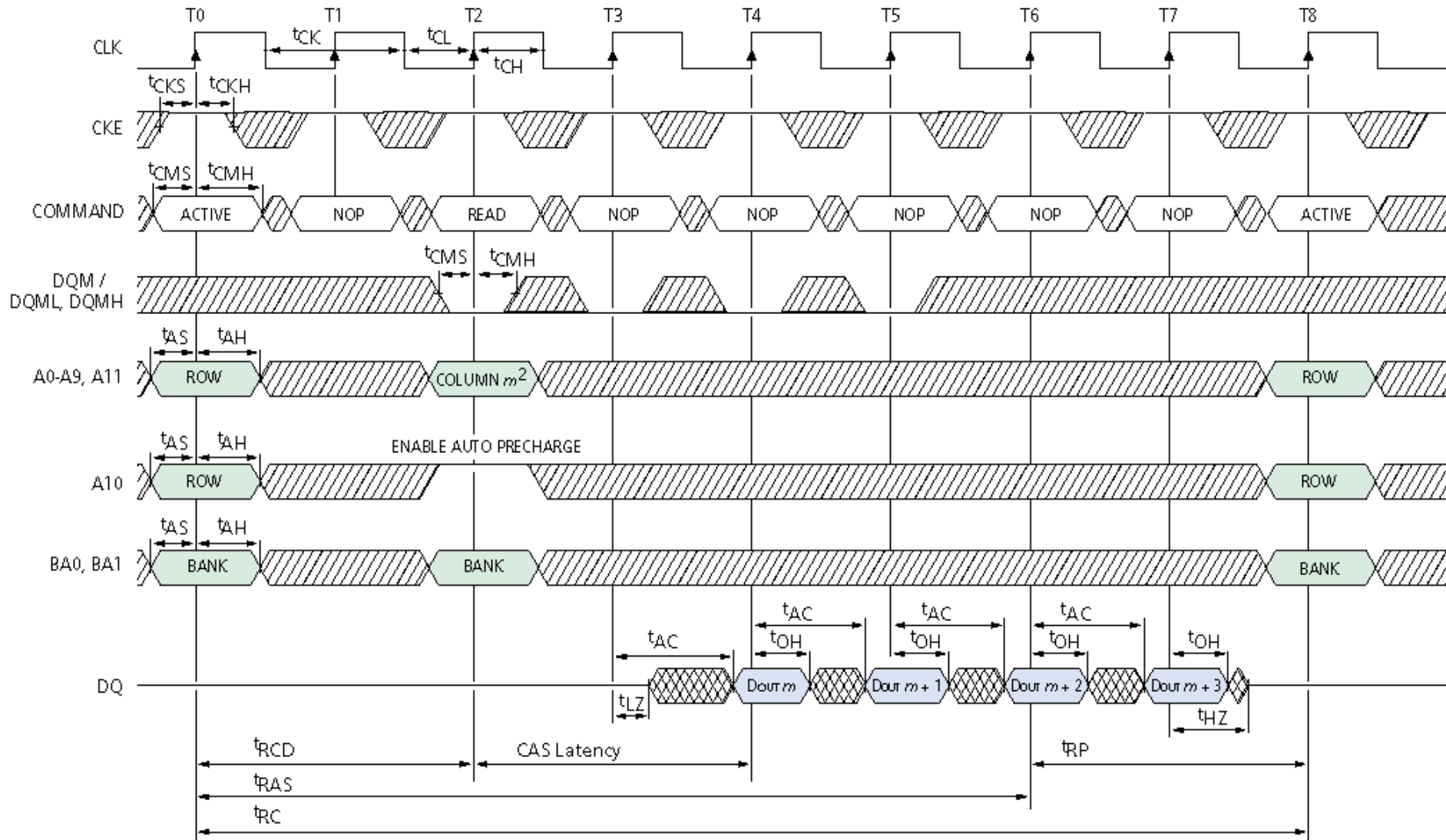
Functional Block Diagram 8 Meg x 16 SDRAM



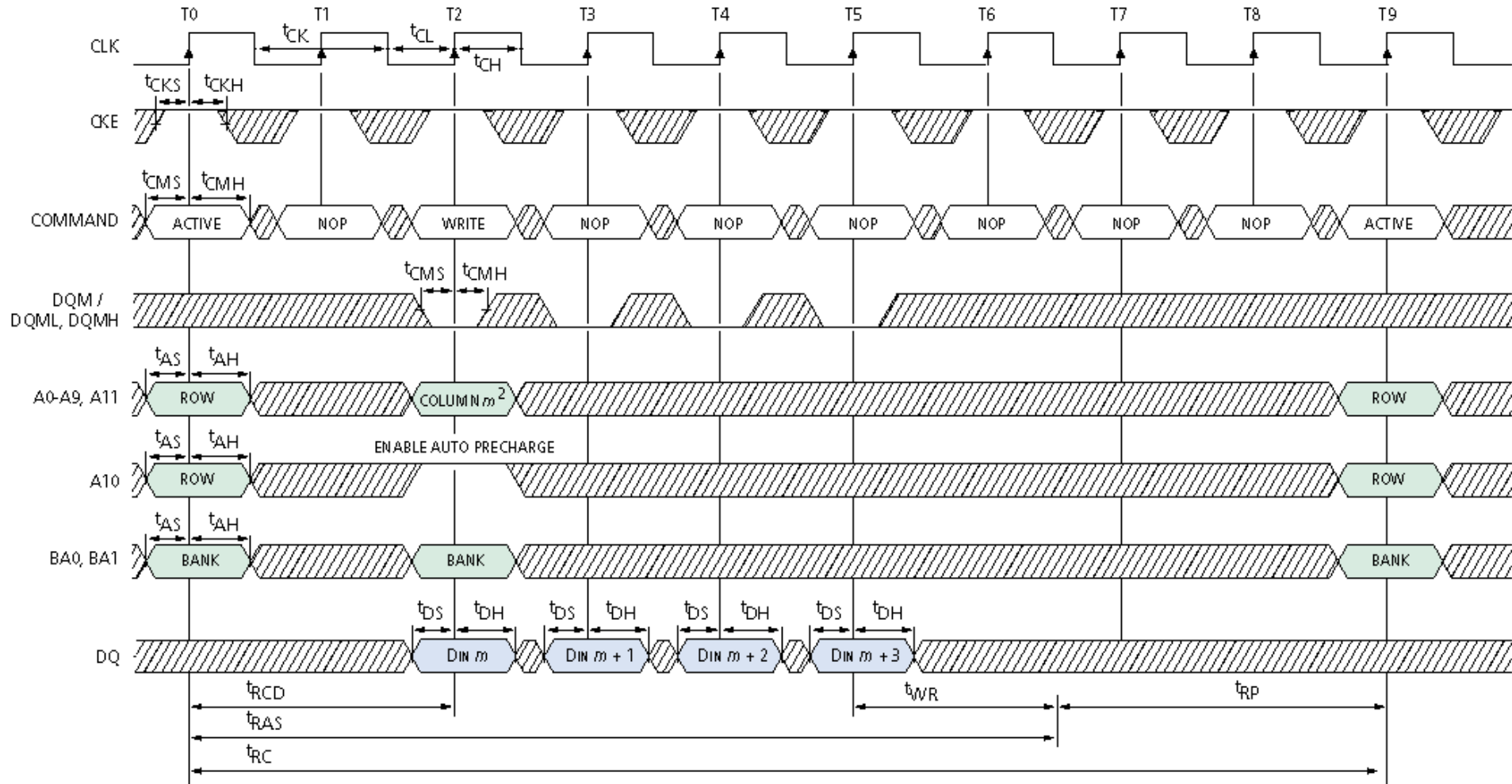
SDRAM Details

- Multiple “banks” of cell arrays are used to reduce access time:
 - Each bank is 4K rows by 512 “columns” by 16 bits (for our part)
- Read and Write operations are split into RAS (row access) followed by CAS (column access)
- These operations are controlled by sending commands
 - Commands are sent using the RAS, CAS, CS, & WE pins.
- Address pins are “time multiplexed”
 - During RAS operation, address lines select the bank and row
 - During CAS operation, address lines select the column.
- “ACTIVE” command “opens” a row for operation
 - transfers the contents of the entire row to a row buffer
- Subsequent “READ” or “WRITE” commands modify the contents of the row buffer.
- For burst reads and writes during “READ” or “WRITE” the starting address of the block is supplied.
 - Burst length is programmable as 1, 2, 4, 8 or a “full page” (entire row) with a burst terminate option.
- Special commands are used for initialization (burst options etc.)
- A burst operation takes $\approx 4 + n$ cycles (for n words)

READ burst (with auto precharge)



WRITE burst (with auto precharge)

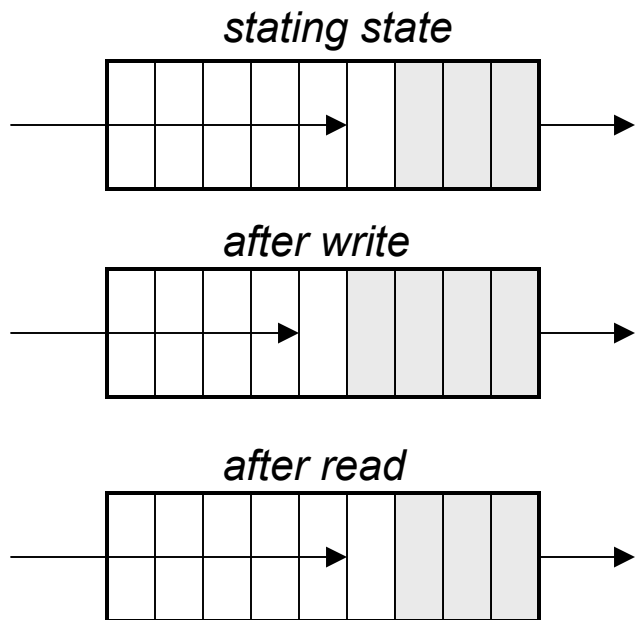


See datasheet for more details.

Verilog simulation models available.

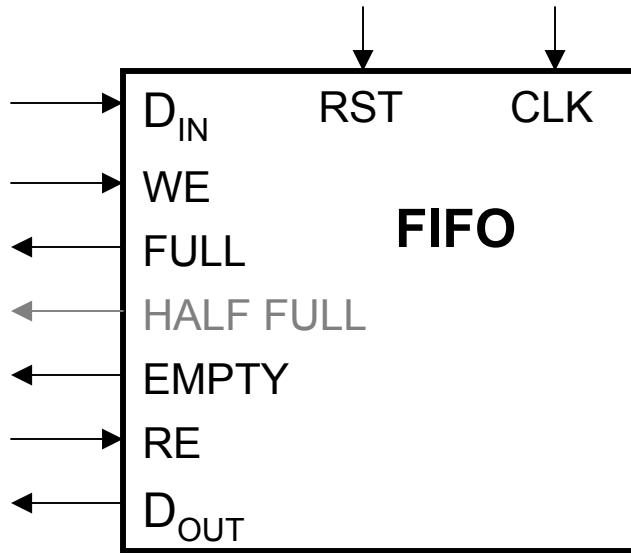
First-in-first-out (FIFO) Memory

- Used to implement *queues*.
- These find common use in computers and communication circuits.
- Generally, used for rate matching data producer and consumer:



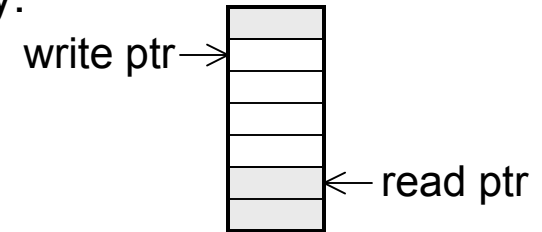
- Producer can perform many writes without consumer performing any reads (or vis versa). However, because of finite buffer size, on average, need equal number of reads and writes.
- Typical uses:
 - interfacing I/O devices. Example network interface. Data bursts from network, then processor bursts to memory buffer (or reads one word at a time from interface). Operations not synchronized.
 - Example: Audio output. Processor produces output samples in bursts (during process swap-in time). Audio DAC clocks it out at constant sample rate.

FIFO Interfaces

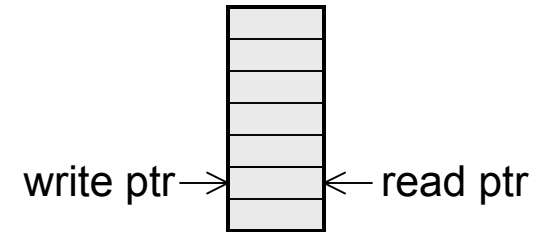


- After write or read operation, $FULL$ and $EMPTY$ indicate status of buffer.
- Used by external logic to control own reading from or writing to the buffer.
- FIFO resets to $EMPTY$ state.
- $HALF FULL$ (or other indicator of partial fullness) is optional.

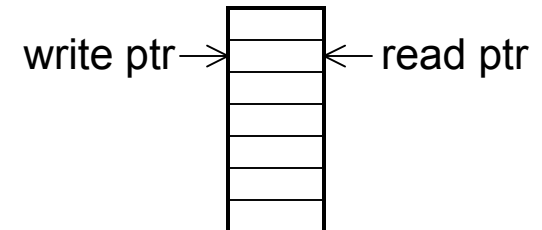
- Address pointers are used internally to keep next write position and next read position into a dual-port memory.



- If pointers equal after write $\Rightarrow FULL$:



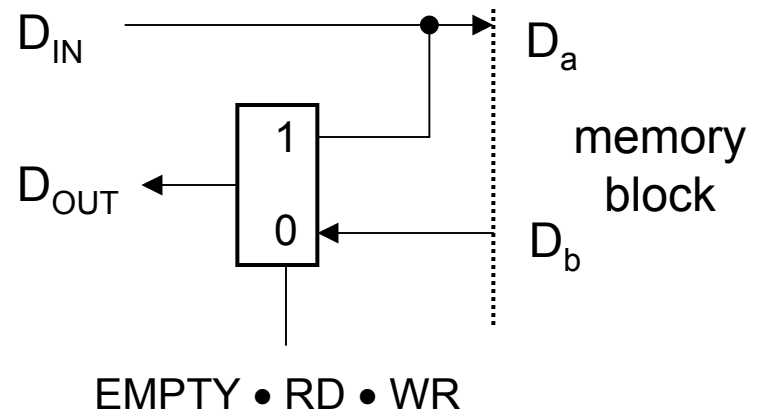
- If pointers equal after read $\Rightarrow EMPTY$:



FIFO Implementation Details

- Assume, dual-port memory with asynchronous read, synchronous write.
- Binary counter for each of read and write address. CEs controlled by WE and RE.
- Equal comparator to see when pointers match.
- Flip-flop each for FULL and EMPTY flags:
- With this memory, read happens before write. Therefore if WE and RD are asserted:
 - When FULL, correct operation will happen.
 - When EMPTY, would like to write before read.
 - Correct this case with a *bypass mux*:

WE	RD	equal	EMPTY _i	FULL _i
0	0	0	0	0
0	0	1	EMPTY _{i-1}	FULL _{i-1}
0	1	0	0	0
0	1	1	1	0
1	0	0	0	0
1	0	1	0	1
1	1	0	0	0
1	1	1	EMPTY _{i-1}	FULL _{i-1}



Xilinx BlockRam Versions

- Our simple version has latency problems.
 - FULL and EMPTY signals are asserted near the end of the clock period.
- Xilinx version solves this by “predicting” when full or empty will be asserted on next write/read operation. Also uses BlockRam with synchronous reads.
- Available on Xilinx website along with “app note” – application note. These are linked to our page.
- Two versions available. Easy to modify to change the width if necessary.
- Will use the “cc” (common clock) version for checkpoint 2.
- Can use the “ic” (independent clock) version later for bridging network to video.